



SMART CONTRACT AUDIT REPORT

for

ONDO



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1 | Introduction

Given the opportunity to review the **Ondo** design document and related smart contract source code, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

1.1 About Ondo

The goal of Ondo is to allow investors to shift the risk and reward balance between each other. In particular, Ondo classifies investors into two groups: the *senior* and *junior* tranche. The *senior* tranche will receive a fixed percentage over their initial investments. The *junior* tranche will receive any excess returns over the senior tranche. This fixed percentage, called the *hurdle rate*, is determined when the product is created. The underlying investment is liquidity provider tokens on decentralized exchanges, e.g. Uniswap, Sushiswap, Balancer, Curve, etc. Liquidity providers inject an equal value of a pair of assets into a liquidity pool. In return they earn fees and incentives for providing liquidity, which is collected by withdrawing liquidity from the pool.

The basic information of Ondo is as follows:

Table 1.1: Basic Information of Ondo

Item	Description
Issuer	Ondo
Type	Ethereum Smart Contract
Platform	Solidity
Audit Method	Whitebox
Latest Audit Report	May 16, 2021

In the following, we show the Git repository of reviewed files and the commit hash value used

in this audit. Note that Ondo assumes a trusted price oracle with timely market price feeds for supported assets and the oracle itself is not part of this audit.

- <https://github.com/ondoprotocol/protocol-audit.git> (29b2c9a)

1.2 About PeckShield

PeckShield Inc. [10] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (<https://t.me/peckshield>), Twitter (<http://twitter.com/peckshield>), or Email (contact@peckshield.com).

Table 1.2: Vulnerability Severity Classification

Impact	High	Critical	High	Medium
	Medium	High	Medium	Low
	Low	Medium	Low	Low
		High	Medium	Low
		Likelihood		

1.3 Methodology

To standardize the evaluation, we define the following terminology based on the OWASP Risk Rating Methodology [9]:

- Likelihood represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

Table 1.3: The Full Audit Checklist

Category	Checklist Items
Basic Coding Bugs	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
	Revert DoS
	Unchecked External Call
	Gasless Send
	Send Instead Of Transfer
	Costly Loop
	(Unsafe) Use Of Untrusted Libraries
	(Unsafe) Use Of Predictable Variables
	Transaction Ordering Dependence
Deprecated Uses	
Semantic Consistency Checks	Semantic Consistency Checks
Advanced DeFi Scrutiny	Business Logics Review
	Functionality Checks
	Authentication Management
	Access Control & Authorization
	Oracle Security
	Digital Asset Escrow
	Kill-Switch Mechanism
	Operation Trails & Event Generation
	ERC20 Idiosyncrasies Handling
	Frontend-Contract Integration
	Deployment Consistency
	Holistic Risk Management
Additional Recommendations	Avoiding Use of Variadic Byte Array
	Using Fixed Compiler Version
	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
Following Other Best Practices	

To evaluate the risk, we go through a checklist of items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- Basic Coding Bugs: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- Semantic Consistency Checks: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [8], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings. Moreover, in case there is an issue that may affect an active protocol that has been deployed, the public version of this report may omit such issue, but will be amended with full details right after the affected protocol is upgraded with respective fixes.

1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary
Configuration	Weaknesses in this category are typically introduced during the configuration of the software.
Data Processing Issues	Weaknesses in this category are typically found in functionality that processes data.
Numeric Errors	Weaknesses in this category are related to improper calculation or conversion of numbers.
Security Features	Weaknesses in this category are concerned with topics like authentication, access control, confidentiality, cryptography, and privilege management. (Software security is not security software.)
Time and State	Weaknesses in this category are related to the improper management of time and state in an environment that supports simultaneous or near-simultaneous computation by multiple systems, processes, or threads.
Error Conditions, Return Values, Status Codes	Weaknesses in this category include weaknesses that occur if a function does not generate the correct return/status code, or if the application does not handle all possible return/status codes that could be generated by a function.
Resource Management	Weaknesses in this category are related to improper management of system resources.
Behavioral Issues	Weaknesses in this category are related to unexpected behaviors from code that an application uses.
Business Logic	Weaknesses in this category identify some of the underlying problems that commonly allow attackers to manipulate the business logic of an application. Errors in business logic can be devastating to an entire application.
Initialization and Cleanup	Weaknesses in this category occur in behaviors that are used for initialization and breakdown.
Arguments and Parameters	Weaknesses in this category are related to improper use of arguments or parameters within function calls.
Expression Issues	Weaknesses in this category are related to incorrectly written expressions within code.
Coding Practices	Weaknesses in this category are related to coding practices that are deemed unsafe and increase the chances that an exploitable vulnerability will be present in the application. They may not directly introduce a vulnerability, but indicate the product has not been carefully developed or maintained.

2 | Findings

2.1 Summary

Here is a summary of our findings after analyzing the implementation of the Ondo protocol. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logic, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings	
Critical	0	
High	1	■
Medium	5	■ ■ ■ ■ ■
Low	5	■ ■ ■ ■ ■
Informational	0	
Total	11	

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in Section 3.

2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 1 high-severity vulnerability, 5 medium-severity vulnerabilities, and 5 low-severity vulnerabilities.

Table 2.1: Key Ondo Audit Findings

ID	Severity	Title	Category	Status
PVE-001	Medium	Improved Logic in Excess Withdrawal	Security Features	Fixed
PVE-002	Low	Duplicate Removal in Registry::tokensDeclaredDead()	Business Logic	Fixed
PVE-003	Low	Explicit Input Validation in AllPair-CCO::transition()	Business Logic	Fixed
PVE-004	Low	Improved Sanity Checks Of System/-Function Parameters	Coding Practices	Fixed
PVE-005	Medium	Incorrect Performance Fee Calculation	Business Logic	Confirmed
PVE-006	Low	Redundant Code Removal	Coding Practices	Fixed
PVE-007	Low	Accommodation of approve() Idiosyncrasies	Coding Practices	Fixed
PVE-008	Medium	Trust Issue of Admin Keys	Security Features	Confirmed
PVE-009	Medium	Possible Front-Running For Reduced Return	Business Logic	Fixed
PVE-010	High	Improved Business Logic in Rollover-CCO::deposit()	Business Logic	Fixed
PVE-011	Medium	Potentially Repeated Excess Returns Of RolloverCCO	Business Logic	Resolved

Beside the identified issues, we emphasize that for any user-facing applications and services, it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms should kick in at the very moment when the contracts are being deployed on mainnet. Please refer to Section 3 for details.

3 | Detailed Results

3.1 Improved Logic in Excess Withdrawal

- ID: PVE-001
- Severity: Medium
- Likelihood: Low
- Impact: High
- Target: BasePairLPStrategy
- Category: Security Features [5]
- CWE subcategory: CWE-287 [2]

Description

The `Ondo` protocol has a number of essential contracts for different functionalities and duties: `strategies`, `AllPairCCO`, `RolloverCCO`, `Registry`, and `TrancheToken`. While examining the `BasePairLPStrategy` contract, we notice one public function `withdrawExcess()` is not properly guarded.

To elaborate, we show below the `withdrawExcess()` routine that is designed to transfer any unused deposits back to the investor. We notice that the current logic employs the `isAuthorized(OLib.CCO_ROLE)` modifier to require `CCO_ROLE` authorization. However, the proper authorization should be more explicit, i.e., `onlyOrigin(_ccoId)`. By doing so, only the origin of the intended `_ccoId` may be able to request the return of excess tokens, not any account with the `CCO_ROLE` authorization.

```
87  /**
88   * @notice Send excess tokens to investor
89   */
90  function withdrawExcess(
91      uint256 _ccold,
92      OLib.Tranche tranche,
93      address to,
94      uint256 amount
95  ) external override isAuthorized(OLib.CCO_ROLE) {
96      CCO storage _cco = ccos[_ccold];
97      if (tranche == OLib.Tranche.Senior) {
98          uint256 excess = _cco.seniorExcess;
99          require(amount <= excess, "Withdrawing too much");
100         _cco.seniorExcess -= amount;
101         _cco.senior.safeTransfer(to, amount);
```

```

102   } else {
103     uint256 excess = _cco.juniorExcess;
104     require(amount <= excess, "Withdrawing too much");
105     _cco.juniorExcess -= amount;
106     _cco.junior.safeTransfer(to, amount);
107   }
108 }

```

Listing 3.1: BasePairLPStrategy::withdrawExcess()

Recommendation Authenticate the caller of `withdrawExcess()` to be `onlyOrigin(_ccoId)`, not the current `isAuthorized(OLib.CCO_ROLE)`.

Status This issue has been fixed in this commit: `1e9db80`.

3.2 Duplicate Removal in Registry::tokensDeclaredDead()

- ID: PVE-002
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: Registry
- Category: Business Logic [7]
- CWE subcategory: CWE-841 [4]

Description

The `undo` protocol has a protocol-wide registry `Registry` with a number of access control-related configurations. Our analysis shows that the contract also maintains an internal list of `TrancheToken` instances that can be recycled for gas efficiency.

To elaborate, we show below the responsible `tokensDeclaredDead()` routine. As the name indicates, this routine adds the given list of tokens into an internal array `deadTokens` where the governance can delete to save gas. However, it comes to our attention that the given list is not validated for any duplicate.

```

88  /**
89   * @notice Manually determine which TrancheToken instances can be recycled
90   * @dev Move into another list where createCCO can delete to save gas. Done manually
91   *       for safety.
92   * @param _tokens List of tokens
93   */
94  function tokensDeclaredDead(address[] calldata _tokens)
95      external
96      onlyGovernance
97  {
98      for (uint256 i = 0; i < _tokens.length; i++) {
99          deadTokens.push(ITrancheToken(_tokens[i]));
100      }

```

100 }

Listing 3.2: Registry :: tokensDeclaredDead()

Recommendation Revise the `tokensDeclaredDead()` logic to not save into the internal array `deadTokens` with duplicates

Status This issue has been fixed in this commit: 2156b55.

3.3 Explicit Input Validation in `AllPairCCO::transition()`

- ID: PVE-003
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: AllPairCCO
- Category: Business Logic [7]
- CWE subcategory: CWE-841 [4]

Description

In the `Ondo` protocol, there is a core `AllPairCCO` contract that contains most of the implementation for `Collateralized Crypto Obligations (CCO)s`. Specifically, rather than creating unique contract instances for each `CCO`, the state for all `CCOs` is stored in a mapping. Each `CCO` has a unique id number created by hashing the metadata on the `CCO`: asset pair, strategy `contract`, start time, investment time, and duration. Each `CCO` has to be one of four self-evident states: `Inactive`, `Deposit`, `Live`, and `Withdraw`. In order to facilitate the (linear) state transition among these four states, a helper routine `transition()` is provided.

To elaborate, we show below the `transition()` helper routine. Our analysis shows one specific transition `Live -> Withdraw` can be improved by explicitly enforcing the required states. In particular, the current logic only enforces `require(curState == OLib.State.Live)` (line 120), which can be improved as `require(curState == OLib.State.Live && _nextState == OLib.State.Withdraw)`.

```

106 // Determine if one can move to a new state. For now the transitions
107 // are strictly linear. No state machines, really.
108 modifier transition(uint256 _ccold, OLib.State _nextState) {
109     CCO storage cco_ = CCOs[_ccold];
110     OLib.State curState = cco_.state;
111     if (_nextState == OLib.State.Live) {
112         require(
113             curState == OLib.State.Deposit,
114             // "Cannot transition to Live from current state"
115             invalidTransitionMsg
116         );
117         require(cco_.investAt <= block.timestamp, "Not yet time to invest");
118     } else {

```

```

119     require(
120         curState == OLib.State.Live,
121         // "Cannot transition to Withdraw from current state"
122         invalidTransitionMsg
123     );
124     require(cco_.redeemAt <= block.timestamp, "Not yet time to redeem");
125 }
126 cco_.state = _nextState;
127 CCOsByState[curState].remove(_ccold);
128 CCOsByState[_nextState].add(_ccold);
129 _;
130 }

```

Listing 3.3: AllPairCCO::transition ()

Recommendation Properly strengthen the `transition()` logic by making state transition explicit.

Status This issue has been fixed in this commit: [7e82cf9](#).

3.4 Improved Sanity Checks For System Parameters

- ID: PVE-004
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: Multiple Contracts
- Category: Coding Practices [6]
- CWE subcategory: CWE-1126 [1]

Description

DeFi protocols typically have a number of system-wide parameters that can be dynamically configured on demand. The `ondo` protocol is no exception. In the following, we examine a number of routines that can be improved to better validate the given input.

The first routine `AllPairCCO::maybeOpenDeposit()` determines whether a given `CCO` can shift to an open state. Besides the current validation, it is better to add the following requirement, i.e., `require(cco_.startAt > 0 && cco_.startAt <= block.timestamp)`, to ensure the `startAt` field is indeed valid.

```

132 // Determine if a CCO can shift to an open state. A CCO is started
133 // in an inactive state. It can only move forward when time has
134 // moved past the starttime.
135 modifier maybeOpenDeposit(uint256 _ccold) {
136     CCO storage cco_ = CCOs[_ccold];
137     if (cco_.state == OLib.State.Inactive) {
138         require(cco_.startAt <= block.timestamp, "Not yet time to enroll");

```

```

139     cco_.state = OLib.State.Deposit;
140     CCOsByState[OLib.State.Inactive].remove(_ccold);
141     CCOsByState[OLib.State.Deposit].add(_ccold);
142   } else if (cco_.state != OLib.State.Deposit) {
143     revert("Invalid operation at current state");
144   }
145   _;
146 }

```

Listing 3.4: AllPairCCO::maybeOpenDeposit()

The second routine `canDeposit()` determines whether the given CCO can accept investor deposits. This routine can be similarly improved by verifying `cco_.startAt > 0 && cco_.startAt <= block.timestamp`.

```

906 function canDeposit(uint256 _ccold) external view override returns (bool) {
907     CCO storage cco_ = CCOs[_ccold];
908     if (cco_.state == OLib.State.Inactive) {
909         return cco_.startAt <= block.timestamp;
910     }
911     return cco_.state == OLib.State.Deposit;
912 }

```

Listing 3.5: AllPairCCO::canDeposit()

The third routine `createAndAddNextCco()` defines the next CCO for this rollover to invest. It is suggested to perform a more thorough validation on the given `_ccoParams` to define a CCO.

```

353 function createAndAddNextCco(
354     uint256 _rolloverId,
355     OLib.CCOParams memory _ccoParams
356 )
357     external
358     noPanic
359     notDead(_rolloverId)
360     onlyCreator(_rolloverId)
361     nonReentrant
362 {
363     Rollover storage rollover_ = rollovers[_rolloverId];
364     _ccoParams.startTime =
365         ccoManager.redeemAt(rollover_.rounds[rollover_.thisRound + 1].ccold) -
366         _ccoParams.enrollment;
367     uint256 newCcold = ccoManager.createCCO(_ccoParams);
368     _addNextCco(_rolloverId, newCcold);
369 }

```

Listing 3.6: RolloverCCO::createAndAddNextCco()

Recommendation Validate any untrusted input before it can be accepted for normal processing. Also, guard any changes on the system-wide parameters to ensure they fall in an appropriate range. Also, consider emitting related events for external monitoring and analytics tools.

Status This issue has been fixed in the following commits: [d1d14e5](#), [8477b11](#), [e802dea](#), and [8f96dbb](#).

3.5 Incorrect Performance Fee Calculation

- ID: PVE-005
- Severity: Low
- Likelihood: Medium
- Impact: Low
- Target: AllPairCCO
- Category: Business Logic [7]
- CWE subcategory: CWE-841 [4]

Description

As mentioned earlier, the `ondo` protocol allows investors to shift the risk and reward balance between each other. Currently, the protocol supports two types of investors: the `senior tranche` and `junior tranche`. The `senior tranche` will receive a fixed percentage over their initial investments. The `junior tranche` will receive any excess returns over the `senior tranche`. This fixed percentage, called the `hurdle rate`, is determined when the product is created.

In particular, we show below the `AllPairCCO::takePerformanceFee()` routine that is used to calculate the performance fee for the strategist. It comes to our attention that the calculated fee is currently based on received amount by `junior tranche`, instead of the earned amount after the `hurdle rate` reduction. As a result, the `senior tranche` may not get the expected `hurdle rate`.

```

789 function takePerformanceFee(CCO storage cco) internal returns (uint256 fee) {
790     fee = 0;
791     if (performanceFeeCollector != address(0)) {
792         Asset storage junior = cco.assets[OLib.Tranche.Junior];
793         uint256 juniorHurdle =
794             junior
795                 .totalInvested
796                 .fromUInt()
797                 .mul((denominator + cco.hurdleRate).fromUInt())
798                 .div(denominator.fromUInt())
799                 .toUInt();
801         if (junior.received > juniorHurdle) {
802             fee = cco
803                 .performanceFee
804                 .fromUInt()
805                 .mul(junior.received.fromUInt())
806                 .div(denominator.fromUInt())
807                 .toUInt();
808             IERC20(junior.token).safeTransferFrom(
809                 address(cco.strategist),
810                 performanceFeeCollector,

```



```

811         fee
812     );
813     }
814 }
815 }

```

Listing 3.7: AllPairCCO::takePerformanceFee()

Recommendation Revise the above calculations to ensure the `senior tranche` can get the expected `hurdle rate`.

Status This issue has been confirmed.

3.6 Redundant Code Removal

- ID: PVE-006
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: Multiple Contracts
- Category: Coding Practices [6]
- CWE subcategory: CWE-563 [3]

Description

The `Undo` protocol makes good use of a number of reference contracts, such as `ERC20`, `SafeERC20`, `SafeMath`, and `AccessControl`, to facilitate its code implementation and organization. For example, the `AllPairCCO` smart contract has so far imported at least five reference contracts. However, we observe the inclusion of certain unused code or the presence of unnecessary redundancies that can be safely removed.

For example, if we examine closely the `UniswapStrategy::constructor()` routine, it contains the repeated initialization of `registry` (line 35).

```

30     constructor(
31         address _registry ,
32         address _router ,
33         address _factory
34     ) BasePairLPStrategy(_registry) {
35         registry = Registry(_registry);
36         uniRouter02 = IUniswapV2Router02(_router);
37         uniFactory = _factory;
38     }

```

Listing 3.8: UniswapStrategy::constructor()

In addition, the `updatePool()` routine can be revised as the requirement on `require(token0 != address(sushiToken) || token1 != address(sushiToken))` is always true – as it is impossible to have a pool with the same `token0` and `token1`.

```

247 function updatePool(address _pool, address[] calldata pathFromSushi)
248     external
249     nonReentrant
250     noPanic
251     isAuthorized(OLib.STRATEGIST_ROLE)
252 {
253     require(pools[_pool]._isSet, "Pool ID not yet registered");
254     address token0 = IUniswapV2Pair(_pool).token0();
255     address token1 = IUniswapV2Pair(_pool).token1();
256     require(
257         token0 != address(sushiToken) & token1 != address(sushiToken),
258         "Should never need to update pool with sushi token"
259     );
260     address endToken = pathFromSushi[pathFromSushi.length - 1];
261     require(
262         IUniswapV2Pair(_pool).token0() == endToken &
263         IUniswapV2Pair(_pool).token1() == endToken,
264         "Not a valid path for pool"
265     );
266     PoolData storage poolData = pools[_pool];
267     delete poolData.pathFromSushi;
268     pools[_pool].pathFromSushi = pathFromSushi;
269 }

```

Listing 3.9: SushiStrategyLP::updatePool()

Similarly, if we examine the `AllPairCCO::depositFromRollover()` routine, the internal requirement `address(cco_.rollover) == msg.sender` (line 441) is redundant as it is already guaranteed by the `onlyRollover(_ccoId, _rolloverId)` modifier.

Recommendation Consider the removal of the redundant code in above routines.

Status This issue has been fixed in the following commits: `1d53499`, `fcc8a89`, and `055bb88`.

3.7 Accommodation of `approve()` Idiosyncrasies

- ID: PVE-007
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: Multiple Contracts
- Category: Coding Practices [6]
- CWE subcategory: CWE-1126 [1]

Description

Though there is a standardized ERC-20 specification, many token contracts may not strictly follow the specification or have additional functionalities beyond the specification. In this section, we examine the `approve()` routine and analyze possible idiosyncrasies from current widely-used token contracts.

In particular, we use the popular stablecoin, i.e., USDT, as our example. We show the related code snippet below. On its entry of `approve()`, there is a requirement, i.e., `require(!((_value != 0) && (allowed[msg.sender][_spender] != 0)))`. This specific requirement essentially indicates the need of reducing the allowance to 0 first (by calling `approve(_spender, 0)`) if it is not, and then calling a second one to set the proper allowance. This requirement is in place to mitigate the known `approve()/transferFrom()` race condition (<https://github.com/ethereum/EIPs/issues/20#issuecomment-263524729>).

```

194  /**
195  * @dev Approve the passed address to spend the specified amount of tokens on behalf
      of msg.sender.
196  * @param _spender The address which will spend the funds.
197  * @param _value The amount of tokens to be spent.
198  */
199  function approve(address _spender, uint _value) public onlyPayloadSize(2 * 32) {

201      // To change the approve amount you first have to reduce the addresses'
202      // allowance to zero by calling 'approve(_spender, 0)' if it is not
203      // already 0 to mitigate the race condition described here:
204      // https://github.com/ethereum/EIPs/issues/20#issuecomment-263524729
205      require(!((_value != 0) && (allowed[msg.sender][_spender] != 0)));

207      allowed[msg.sender][_spender] = _value;
208      Approval(msg.sender, _spender, _value);
209  }

```

Listing 3.10: USDT Token **Contract**

Because of that, a normal call to `approve()` with a currently non-zero allowance may fail. In the following, we use the `SushiStrategyLP::addLiquidity()` routine as an example. This routine is designed to approve a specific token for swap contract. To accommodate the specific idiosyncrasy, for each `safeIncreaseAllowance()`, there is a need to `approve()` twice (lines 445 – 446): the first one reduces the allowance to 0; and the second one sets the new allowance.

```

440  function addLiquidity(
441      address token0,
442      address token1,
443      uint256 amt0,
444      uint256 amt1,
445      uint256 minOut0,
446      uint256 minOut1
447  )
448  internal
449  returns (
450      uint256 out0,
451      uint256 out1,
452      uint256 lp
453  )
454  {
455      IERC20(token0).safeIncreaseAllowance(address(sushiRouter), amt0);
456      IERC20(token1).safeIncreaseAllowance(address(sushiRouter), amt1);

```

```

457     (out0 , out1 , lp) = sushiRouter.addLiquidity(
458         token0 ,
459         token1 ,
460         amt0 ,
461         amt1 ,
462         minOut0 ,
463         minOut1 ,
464         address(this) ,
465         block.timestamp
466     );
467 }

```

Listing 3.11: SushiStrategyLP::addLiquidity()

Moreover, it is important to note that for certain non-compliant ERC20 tokens (e.g., USDT), the `transfer()` function does not have a return value. However, the `IERC20` interface has defined the `transfer()` interface with a `bool` return value. As a result, the call to `transfer()` may expect a return value. With the lack of return value of USDT's `transfer()`, the call will be unfortunately reverted.

Because of that, a normal call to `transfer()` is suggested to use the safe version, i.e., `safeTransfer()`. In essence, it is a wrapper around ERC20 operations that may either throw on failure or return false without reverts. Moreover, the safe version also supports tokens that return no value (and instead revert or throw on failure). Note that non-reverting calls are assumed to be successful. To use this library you can add a using `SafeERC20` for `IERC20`. Similarly, there is a safe version of `approve()/transferFrom()` as well, i.e., `safeApprove()/safeTransferFrom()`.

Recommendation Accommodate the above-mentioned idiosyncrasy about ERC20-related `approve()/transfer()/transferFrom()`.

Status This issue has been fixed in this commit: 074dd9b.

3.8 Trust Issue of Admin Keys

- ID: PVE-008
- Severity: Medium
- Likelihood: Medium
- Impact: Medium
- Target: Multiple Contracts
- Category: Security Features [5]
- CWE subcategory: CWE-287 [2]

Description

In the `ondo` protocol, there is a special administrative account `gov` with `GOVERNANCE_ROLE`. This `gov` account plays a critical role in governing and regulating the system-wide operations (e.g., assign various roles, specify the swap path, and set rollovers). It also has the privilege to regulate or

govern the flow of assets among the involved components in the protocol. And the presence of an administrative account can allow for emergency operations.

We emphasize that current privilege assignment is necessary and required for proper protocol operation. However, it is worrisome if the `gov` is not governed by a DAO-like structure. The discussion with the team has confirmed that the `gov` will be managed by a multi-sig account. Note that a compromised `gov` account is capable of modifying current protocol configuration with adverse consequences on user funds.

Recommendation Promptly transfer the `gov` privilege to the intended DAO-like governance contract.

Status This issue has been confirmed.

3.9 Possible Front-Running For Reduced Return

- ID: PVE-009
- Severity: Medium
- Likelihood: Medium
- Impact: Medium
- Target: Multiple Contracts
- Category: Business Logic [7]
- CWE subcategory: CWE-841 [4]

Description

In the `Ondo` protocol, there is a `SushiStrategyLP` strategy that has an additional method `harvest()` to occasionally convert earned `Sushi` into a balance of `senior` and `junior` assets to reinvest into LP tokens, which are then placed into `Masterchef` to earn `Sushi`.

To elaborate, we show below the `SushiStrategyLP::harvest()` routine. This routine delegates the call to an internal `_compound()` handler to reinvest `sushi/xsushi` into LP tokens.

```

352  /**
353   * @notice Periodically reinvest sushi/xsushi into LP tokens
354   * @param pool Sushiswap pool to reinvest
355   */
356  function harvest(address pool) external isAuthorized(OLib.STRATEGIST_ROLE) {
357      PoolData storage poolData = pools[pool];
358      _compound(IERC20(pool), poolData);
359  }
```

Listing 3.12: `SushiStrategyLP::harvest()`

```

277  function _compound(IERC20 pool, PoolData storage poolData) internal {
278      uint256 sushiAmt = sushiToken.balanceOf(address(this));
279      masterChef.deposit(poolData.pid, 0); // Called to trigger update in amount of sushi
      truly available now
```

```

280 xSushi.leave(poolData.pendingXSushi);
281 poolData.pendingXSushi = 0;
282 sushiAmt = sushiToken.balanceOf(address(this)) - sushiAmt;
283 if (sushiAmt > 0) {
284     address[] memory pathFromSushi = getSushiPath(poolData.pathFromSushi);
285     address tokenA = pathFromSushi[pathFromSushi.length - 1];
286     address tokenB = IUniswapV2Pair(address(pool)).token0();
287     if (tokenB == tokenA) tokenB = IUniswapV2Pair(address(pool)).token1();
288     uint256 amt0;
289     if (tokenA == address(sushiToken)) {
290         amt0 = sushiAmt;
291     } else {
292         amt0 = swapExactIn(sushiAmt, 0, pathFromSushi);
293     }
294     uint256 amt0ToSwap;
295     (uint256 reserves0, ) =
296         SushiSwapLibrary.getReserves(sushiFactory, tokenA, tokenB);
297     amt0 == (amt0ToSwap = calculateSwapInAmount(reserves0, amt0));
298     uint256 amt1 = swapExactIn(amt0ToSwap, 0, getPath(tokenA, tokenB));
299     (, , uint256 lpAmt) = addLiquidity(tokenA, tokenB, amt0, amt1, 0, 0);
300     // TODO: do something with excess - will be extremely minimal though (<2)
301     poolData.totalLp += lpAmt;
302 }
303 pool.safeIncreaseAllowance(
304     address(masterChef),
305     pool.balanceOf(address(this))
306 );
307 masterChef.deposit(poolData.pid, pool.balanceOf(address(this)));
308 }

```

Listing 3.13: SushiStrategyLP::_compound()

We notice the conversion is routed to `SushiSwap` in order to swap one token to another for liquidity addition. And the swap operation does not specify any restriction on possible slippage and is therefore vulnerable to possible front-running attacks, resulting in a smaller gain for this round of yielding.

Note that this is a common issue plaguing current AMM-based DEX solutions. Specifically, a large trade may be sandwiched by a preceding sell to reduce the market price, and a tailgating buy-back of the same amount plus the trade amount. Such sandwiching behavior unfortunately causes a loss and brings a smaller return as expected to the trading user or the `strategy` contract in our case because the swap rate is lowered by the preceding sell. As a mitigation, we may consider specifying the restriction on possible slippage caused by the trade or referencing the TWAP or time-weighted average price of `UniswapV2`. Nevertheless, we need to acknowledge that this is largely inherent to current blockchain infrastructure and there is still a need to continue the search efforts for an effective defense.

Recommendation Develop an effective mitigation to the above front-running attack to better protect the interests of farming users.

Status This issue has been fixed in this commit: 6b7f856.

3.10 Improved Business Logic in RolloverCCO::deposit()

- ID: PVE-010
- Severity: High
- Likelihood: Medium
- Impact: High
- Target: RolloverCCO
- Category: Business Logic [7]
- CWE subcategory: CWE-841 [4]

Description

The `undo` protocol also supports `RolloverCCO` that can automate the investment process by investing in a series of similar CCOs. Investors can deposit money into the `RolloverCCO`, get a token to represent their stake, and redeem it later to collect their investment plus profits (or losses). In the following, we examine the deposit logic.

To elaborate, we use `RolloverCCO::deposit()` routine. The logic is rather straightforward in allowing investors to deposit tokens into a queue to get invested in the next CCO. However, it comes to our attention that when there is an excess and the excess is more than the deposited amount for investment, there is a need to reset `_amount = 0`. The reset operation is missing in current logic, resulting in possible loss for the depositing user.

```

437 function deposit(
438     uint256 _rolloverId ,
439     OLib.Tranche _tranche ,
440     uint256 _amount
441 ) external noPanic notDead(_rolloverId) nonReentrant {
442     Rollover storage rollover_ = rollovers[_rolloverId];
443     if (rollover_.investorLastUpdates[_tranche][msg.sender] == 0) {
444         rollover_.investorLastUpdates[_tranche][msg.sender] = 1;
445     }
446     {
447         Round storage round_ = rollover_.rounds[rollover_.thisRound + 1];
448         uint256 ccold = round_.ccold;
449         require(ccold != 0, "No CCO to deposit in yet");
450         require(ccoManager.canDeposit(ccold), "CCO not in deposit state");
451         TrancheRound storage trancheRound_ = round_.tranches[_tranche];
452         OLib.Investor storage investor_ = trancheRound_.investors[msg.sender];
453         uint256 total = trancheRound_.newDeposited += _amount;
454         uint256 userSum =
455             investor_.userSums.length > 0
456             ? investor_.userSums[investor_.userSums.length - 1] + _amount
457             : _amount;
458         investor_.prefixSums.push(total);
459         investor_.userSums.push(userSum);

```

```

460 }
461 if (
462     rollover_.investorLastUpdates[_tranche][msg.sender] <
463     rollover_.thisRound + 1
464 ) {
465     (uint256 shares, uint256 excess) =
466         _updateInvestor(msg.sender, _rolloverId, _tranche);
467     rollover_.investorLastUpdates[_tranche][msg.sender] =
468         rollover_.thisRound +
469         1;
470     if (excess > _amount) {
471         rollover_.assets[_tranche].safeTransfer(msg.sender, excess - _amount);
472     } else {
473         _amount -= excess;
474     }
475     if (shares > 0) {
476         rollover_.rolloverTokens[_tranche].mint(msg.sender, shares);
477     }
478 }
479 rollover_.assets[_tranche].safeTransferFrom(
480     msg.sender,
481     address(this),
482     _amount
483 );
484 emit Deposited(msg.sender, _rolloverId, _amount);
485 }

```

Listing 3.14: RolloverCCO::deposit()

Recommendation Revise the `deposit()` logic in `RolloverCCO` to transfer user funds (via `transferFrom()` at line 479) when the excess is already more than the intended amount for deposit.

Status This issue has been fixed in this commit: 5162589.

3.11 Potentially Repeated Excess Returns Of RolloverCCO

- ID: PVE-011
- Severity: Medium
- Likelihood: Medium
- Impact: Medium
- Target: `RolloverCCO`
- Category: Business Logic [7]
- CWE subcategory: CWE-841 [4]

Description

As mentioned in Section 3.10, the `Ono` protocol supports `RolloverCCO` that can automate the investment process by investing in a series of similar CCOs. Our analysis shows an issue that may result in multiple returns of the same excess amount.

To elaborate, the protocol starts with the round number 0 and participating users can call `deposit()` to invest funds into the `RolloverCCO` contract. When the `migrate()` function is called to initiate the investment into the next `CCO`, the first-time investment executes the `_firstInvest()` routine (line 618), which properly advances the `rollover_.thisRound` to 1 (line 771).

```

611  function migrate(uint256 _rolloverId, SlippageSettings memory _slippage)
612      external
613      noPanic
614      notDead(_rolloverId)
615      onlyCreator(_rolloverId)
616  {
617      if (rollovers[_rolloverId].thisRound == 0) {
618          _firstInvest(
619              _rolloverId,
620              _slippage.seniorMinInvest,
621              _slippage.juniorMinInvest
622          );
623      } else {
624          _migrate(_rolloverId, _slippage);
625      }
626  }

```

Listing 3.15: `RolloverCCO::migrate()`

```

735  function _firstInvest(
736      uint256 _rolloverId,
737      uint256 _seniorMinInvest,
738      uint256 _juniorMinInvest
739  ) internal {
740      Rollover storage rollover_ = rollovers[_rolloverId];
741      Round storage round_ = rollover_.rounds[1];
742      uint256 ccold = round_.ccold;
743      require(round_.ccold != 0, "CCO not set");
744      TrancheRound storage srRound = round_.tranches[OLib.Tranche.Senior];
745      TrancheRound storage jrRound = round_.tranches[OLib.Tranche.Junior];
746      rollover_.assets[OLib.Tranche.Senior].safeIncreaseAllowance(
747          address(ccoManager),
748          srRound.newDeposited
749      );
750      rollover_.assets[OLib.Tranche.Junior].safeIncreaseAllowance(
751          address(ccoManager),
752          jrRound.newDeposited
753      );
754      ccoManager.depositFromRollover(
755          ccold,
756          _rolloverId,
757          srRound.newDeposited,
758          jrRound.newDeposited
759      );
760      ccoManager.invest(ccold, _seniorMinInvest, _juniorMinInvest);
761      (srRound.invested, jrRound.invested) = ccoManager.rolloverClaim(
762          ccold,

```

```

763     _rolloverId
764 );
765 srRound.deposited = srRound.invested;
766 jrRound.deposited = jrRound.invested;
767 srRound.newInvested = srRound.invested;
768 jrRound.newInvested = jrRound.invested;
769 srRound.shares = srRound.invested;
770 jrRound.shares = jrRound.invested;
771 rollover_.thisRound = 1;
772 }

```

Listing 3.16: RolloverCCO::_firstInvest ()

After that, a user may call `claim()`, which cascadingly calls `updateInvestorDistribute()` to claim the excess amount (line 541). Moreover, if the user calls `deposit()` one more time, the `if` condition (lines 461 – 463) is satisfied to execute the `then`-branch, which includes the `_updateInvestor()` execution. This execution allows the current depositing user to claim the excess amount a second time.

```

492 function claim(uint256 _rolloverId, OLib.Tranche _tranche)
493     external
494     noPanic
495     notDead(_rolloverId)
496     nonReentrant
497 {
498     Rollover storage rollover_ = rollovers[_rolloverId];
499     if (
500         rollover_.investorLastUpdates[_tranche][msg.sender] !=
501         rollover_.thisRound + 1
502     ) {
503         _updateInvestorDistribute(msg.sender, _rolloverId, _tranche);
504     }
505 }

```

Listing 3.17: RolloverCCO::claim()

```

437 function deposit(
438     uint256 _rolloverId,
439     OLib.Tranche _tranche,
440     uint256 _amount
441 ) external noPanic notDead(_rolloverId) nonReentrant {
442     Rollover storage rollover_ = rollovers[_rolloverId];
443     if (rollover_.investorLastUpdates[_tranche][msg.sender] == 0) {
444         rollover_.investorLastUpdates[_tranche][msg.sender] = 1;
445     }
446     {
447         Round storage round_ = rollover_.rounds[rollover_.thisRound + 1];
448         uint256 ccold = round_.ccold;
449         require(ccold != 0, "No CCO to deposit in yet");
450         require(ccoManager.canDeposit(ccold), "CCO not in deposit state");
451         TrancheRound storage trancheRound_ = round_.tranches[_tranche];

```

```
452     OLib.Investor storage investor_ = trancheRound_.investors[msg.sender];
453     uint256 total = trancheRound_.newDeposited += _amount;
454     uint256 userSum =
455         investor_.userSums.length > 0
456         ? investor_.userSums[investor_.userSums.length - 1] + _amount
457         : _amount;
458     investor_.prefixSums.push(total);
459     investor_.userSums.push(userSum);
460 }
461 if (
462     rollover_.investorLastUpdates[_tranche][msg.sender] <
463     rollover_.thisRound + 1
464 ) {
465     (uint256 shares, uint256 excess) =
466         _updateInvestor(msg.sender, _rolloverId, _tranche);
467     rollover_.investorLastUpdates[_tranche][msg.sender] =
468         rollover_.thisRound +
469         1;
470     if (excess > _amount) {
471         rollover_.assets[_tranche].safeTransfer(msg.sender, excess - _amount);
472     } else {
473         _amount -= excess;
474     }
475     if (shares > 0) {
476         rollover_.rolloverTokens[_tranche].mint(msg.sender, shares);
477     }
478 }
479 rollover_.assets[_tranche].safeTransferFrom(
480     msg.sender,
481     address(this),
482     _amount
483 );
484 emit Deposited(msg.sender, _rolloverId, _amount);
485 }
```

Listing 3.18: RolloverCCO::deposit()

Recommendation Revise the `deposit()` to avoid repeated claims of any excess amount after investment.

Status This issue has been resolved.

4 | Conclusion

In this audit, we have analyzed the Ondo design and implementation. The system presents a unique, robust offering as a decentralized protocol to allow investors to shift the risk and reward balance between each other. The current code base is well structured and neatly organized. Those identified issues are promptly confirmed and fixed.

Meanwhile, we need to emphasize that [Solidity](#)-based smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.



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